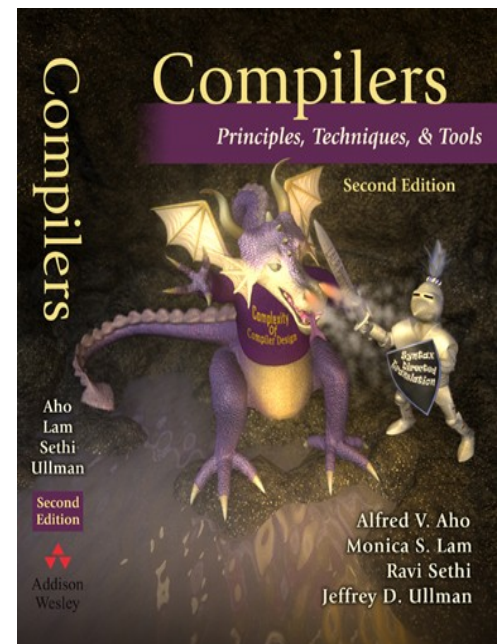


Compiler Design (40-414)

- Main Text Book:
Compilers: Principles, Techniques & Tools, 2nd ed.,
Aho, Lam, Sethi, and Ullman, 2007
- Evaluation:
 - Midterm Exam 35%
 - Final Exam 35%
 - Assignments 10%
 - Project 20%



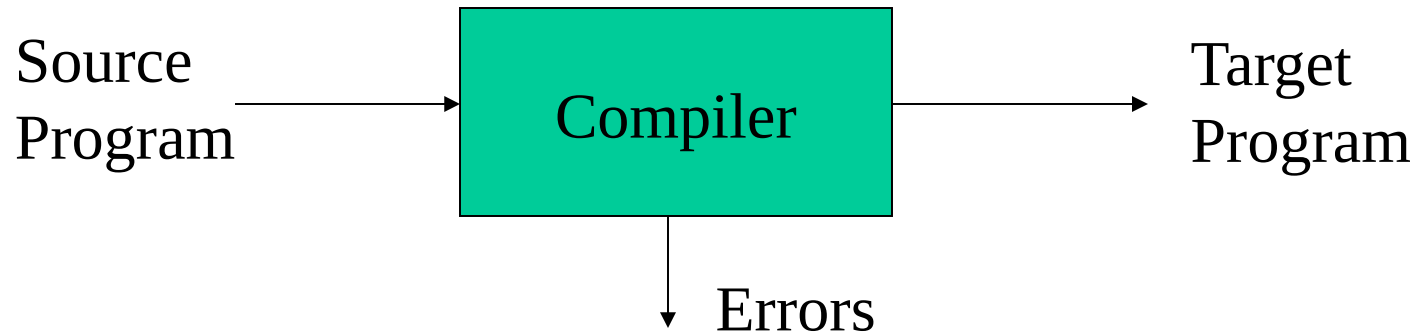
Compiler learning

- Isn't it an old discipline?
 - Yes, it is a well-established discipline
 - Algorithms, methods and techniques were developed in early stages of computer science
 - There are many compilers around, and
 - many tools to generate them automatically
- So, why we need to learn it?
 - Although you may never write a full compiler
 - But the techniques we learn is useful in many tasks like:
 - writing an interpreter for a scripting language,
 - validation checking for forms, and
 - so on

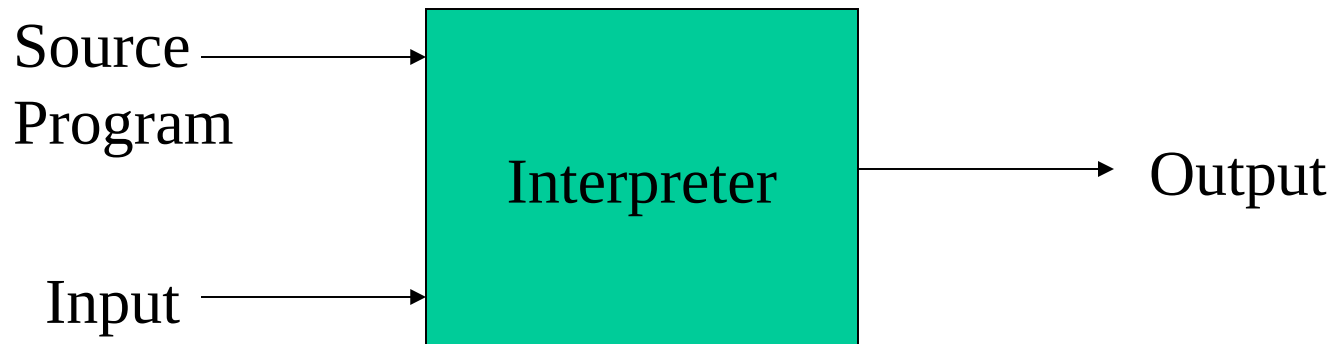
Terminology

- Compiler:
 - a program that translates an *executable* program in a *source language* (usually high level) into an equivalent *executable* program in a *target language* (usually low level)
- Interpreter:
 - a program that reads an *executable* program and produces the results of running that program
 - usually, this involves executing the source program in some fashion
- Our course is mainly about compilers but many of the same issues arise in interpreters

A Compiler

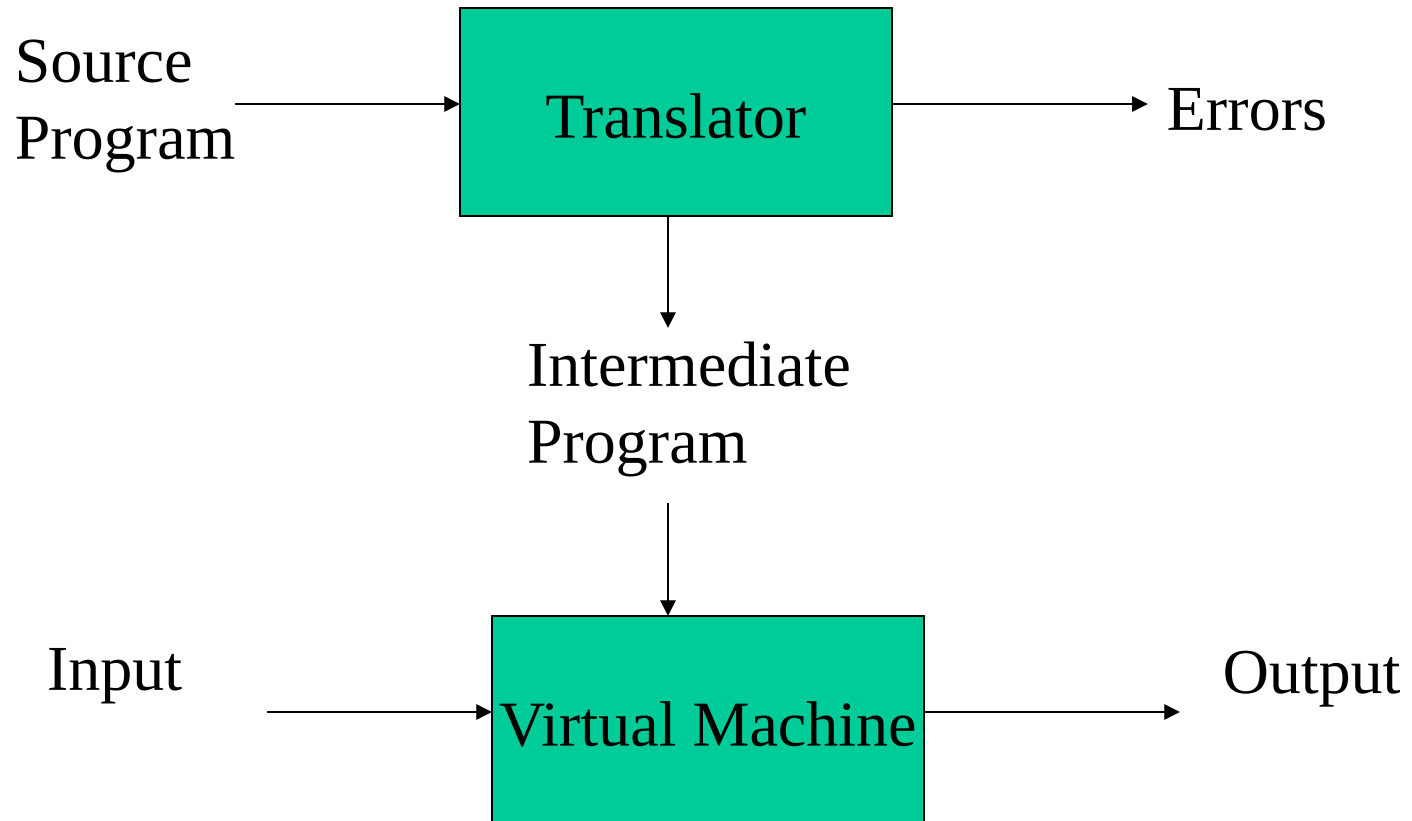


An Interpreter



- Translates line by line
- Executes each translated line immediately
- Execution is slower because translation is repeated
- But, usually give better error diagnostics than a compiler

A Hybrid Compiler



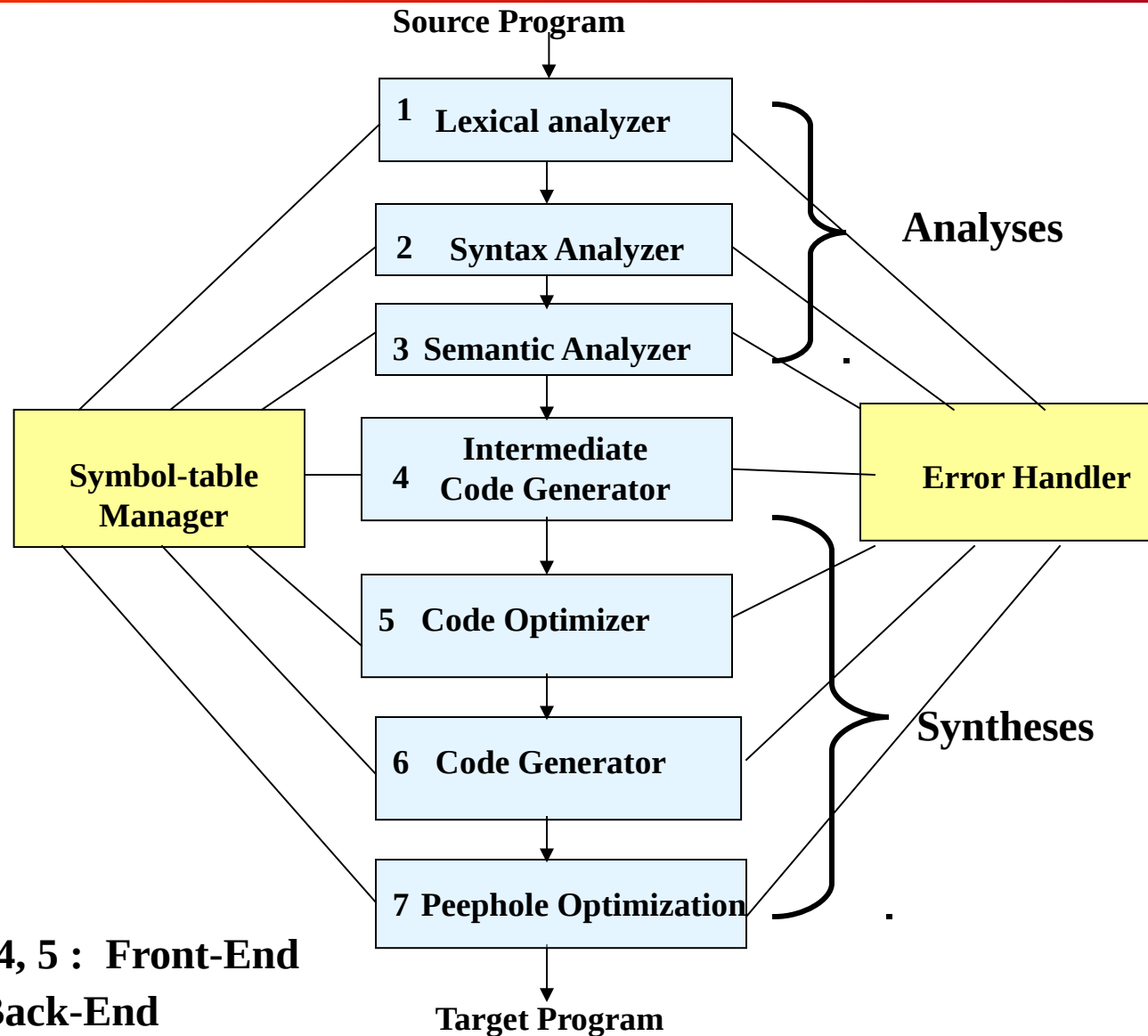
Classifications of Compilers

- There are different types of Compilers:

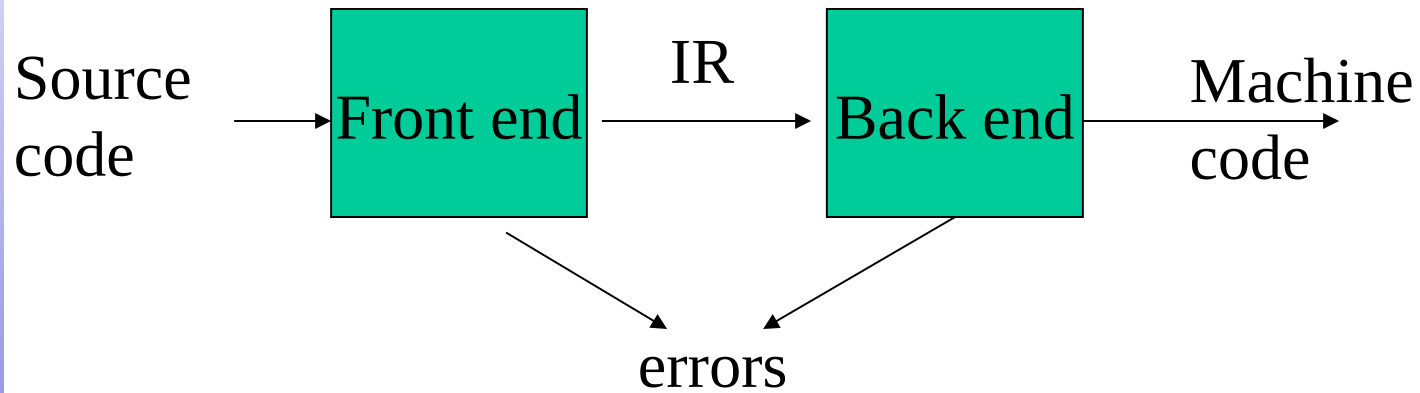
Single Pass } **Construction**
Multiple Pass }

Absolute (e.g., *.com) } **Type of**
Relocateable (e.g., *.exe) } **produced code**

The Many **Phases** of a Compiler

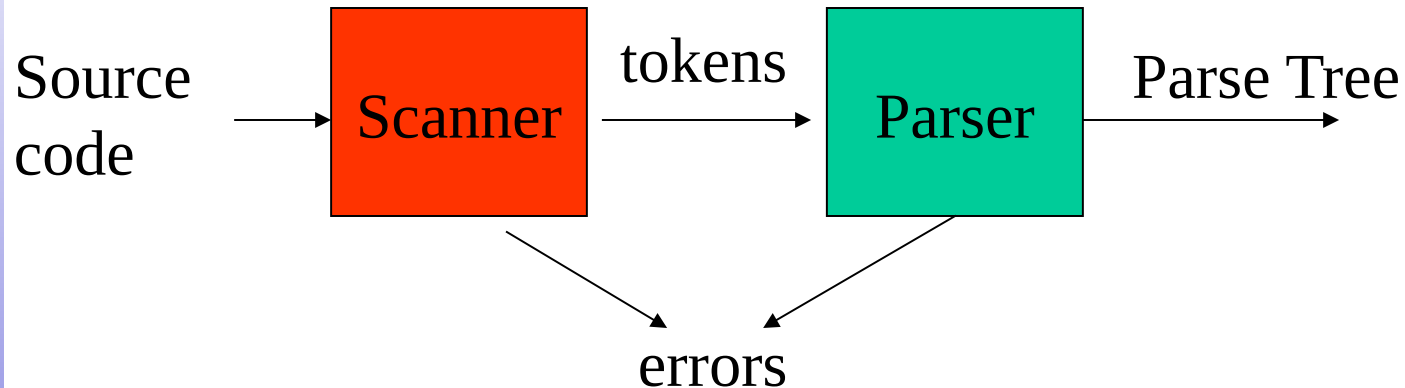


Front-end, Back-end division



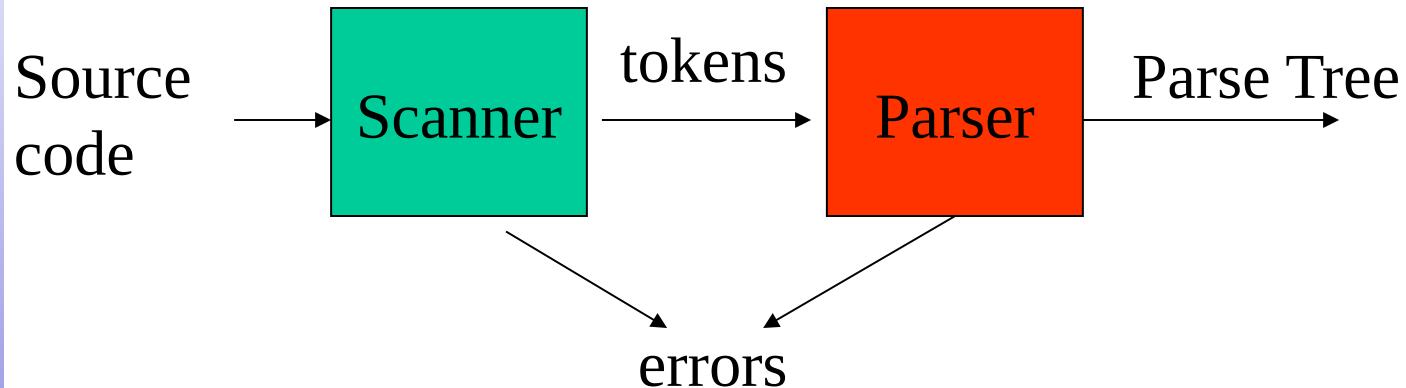
- Front end maps legal code into IR
- Back end maps IR onto target machine
- Simplifies retargeting
- Allows multiple front ends

Front end



- Scanner:
 - ❑ Maps characters into tokens – the basic unit of syntax
 - $x = x + y$ becomes $\langle \text{id}, x \rangle = \langle \text{id}, x \rangle + \langle \text{id}, y \rangle$
 - ❑ Typical tokens: number, id, +, -, *, /, do, end
 - ❑ Eliminate white space (tabs, blanks, comments)
- A key issue is speed so instead of using a tool like LEX it sometimes needed to write your own scanner

Front end



Parser:

- ❑ Recognize context-free syntax
- ❑ Guide context-sensitive analysis
- ❑ Construct IR
- ❑ Produce meaningful error messages
- ❑ Attempt error correction

There are parser generators like YACC which automates much of the work

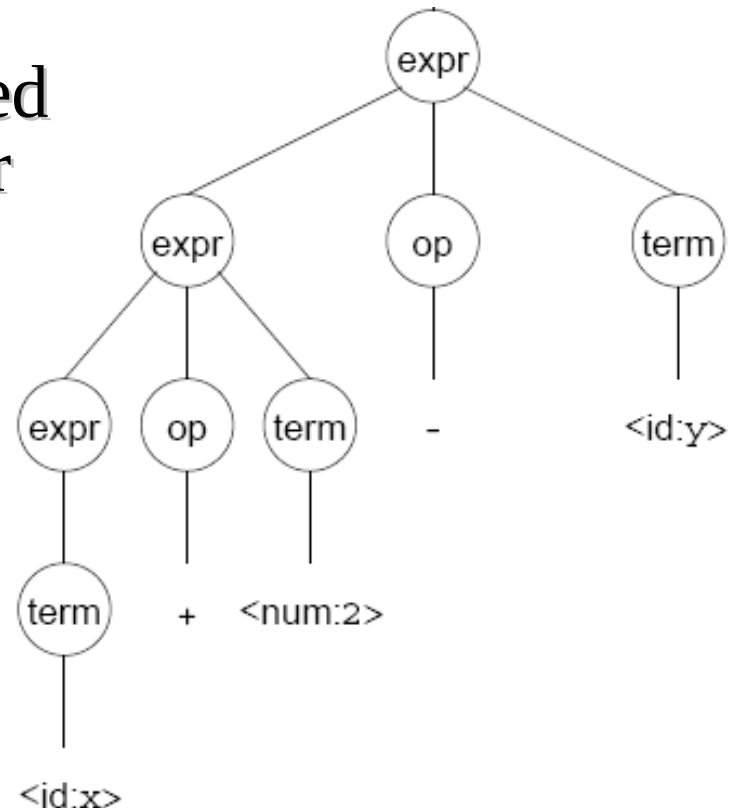
Front end

- Context free grammars are used to represent programming language syntaxes:

$$\langle \text{expr} \rangle ::= \langle \text{expr} \rangle \langle \text{op} \rangle \langle \text{term} \rangle \mid \langle \text{term} \rangle$$
$$\langle \text{term} \rangle ::= \langle \text{number} \rangle \mid \langle \text{id} \rangle$$
$$\langle \text{op} \rangle ::= + \mid -$$

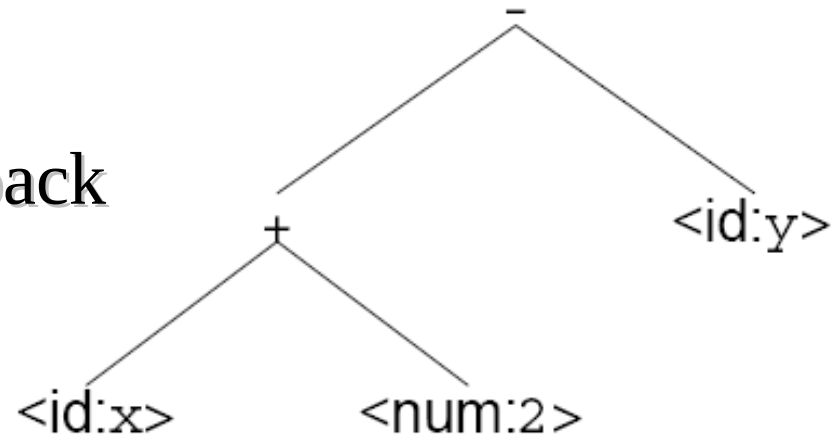
Front end

- A parser tries to map a program to the syntactic elements defined in the grammar
- A parse can be represented by a tree called a parse or syntax tree

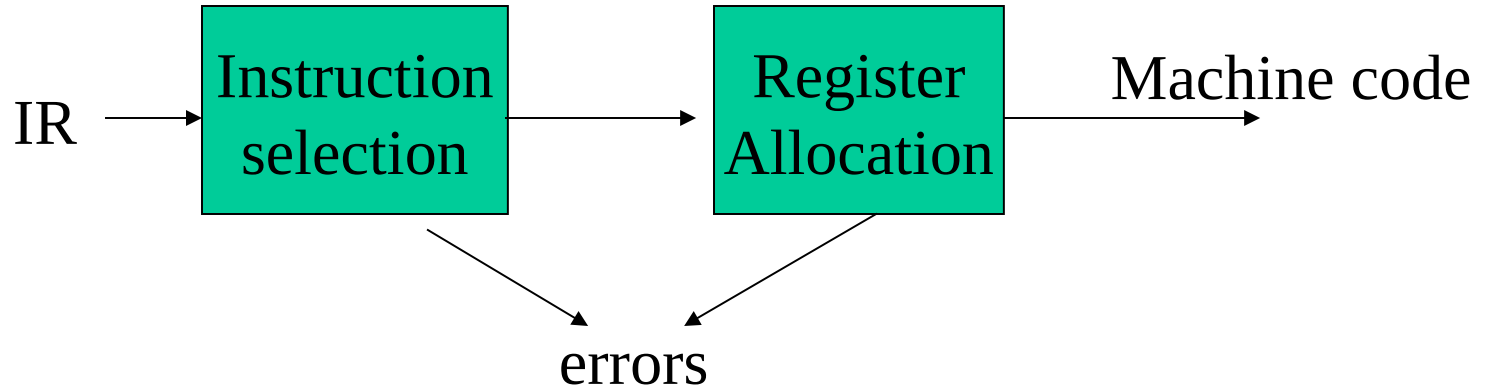


Front end

- A parse tree can be represented more compactly referred to as Abstract Syntax Tree (AST)
- AST can be used as IR between front end and back end

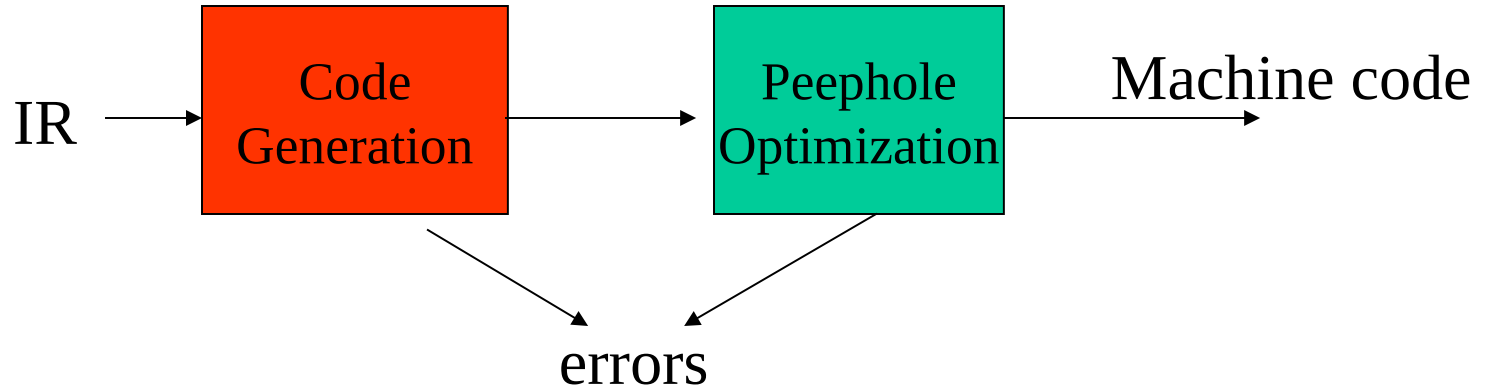


Back end



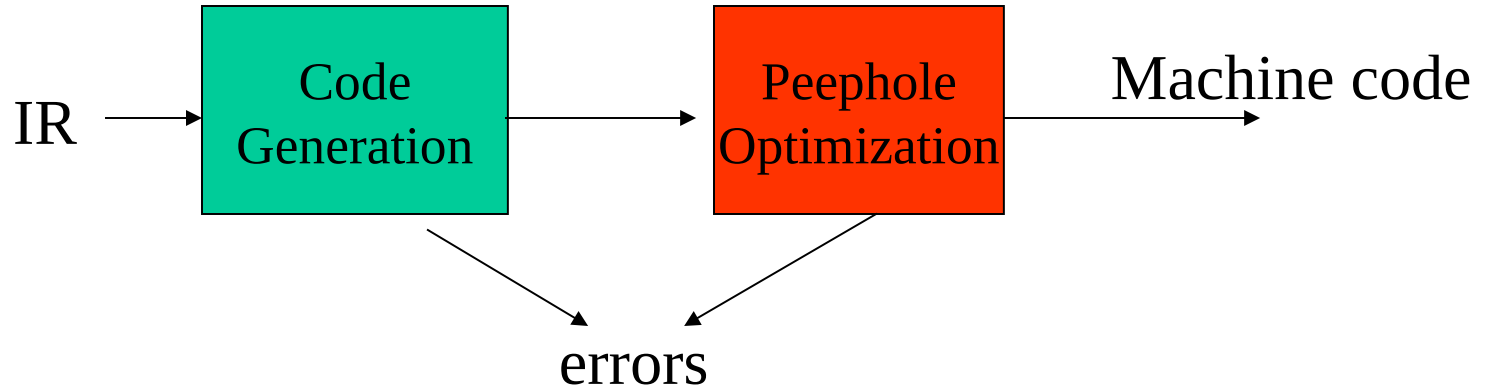
- Translate IR into target machine code
- Choose instructions for each IR operation
- Decide what to keep in registers at each point

Back end



- Produce compact fast code
- Use available addressing modes

Back end



- Limited resources
- Optimal allocation is difficult

The Analysis Task For Compilation

- Three Phases:

- Lexical Analysis:

- Left-to-right Scan to Identify Tokens
 - token: sequence of chars having a collective meaning

- Syntax Analysis:

- Grouping of Tokens Into Meaningful Collection

- Semantic Analysis:

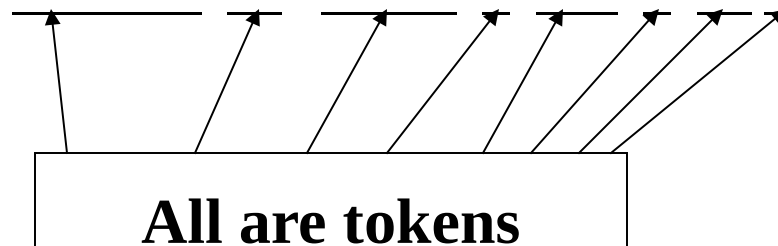
- Checking to ensure Correctness of Components

Phase 1. **Lexical Analysis**

Easiest Analysis - Identify tokens which are the basic building blocks

For

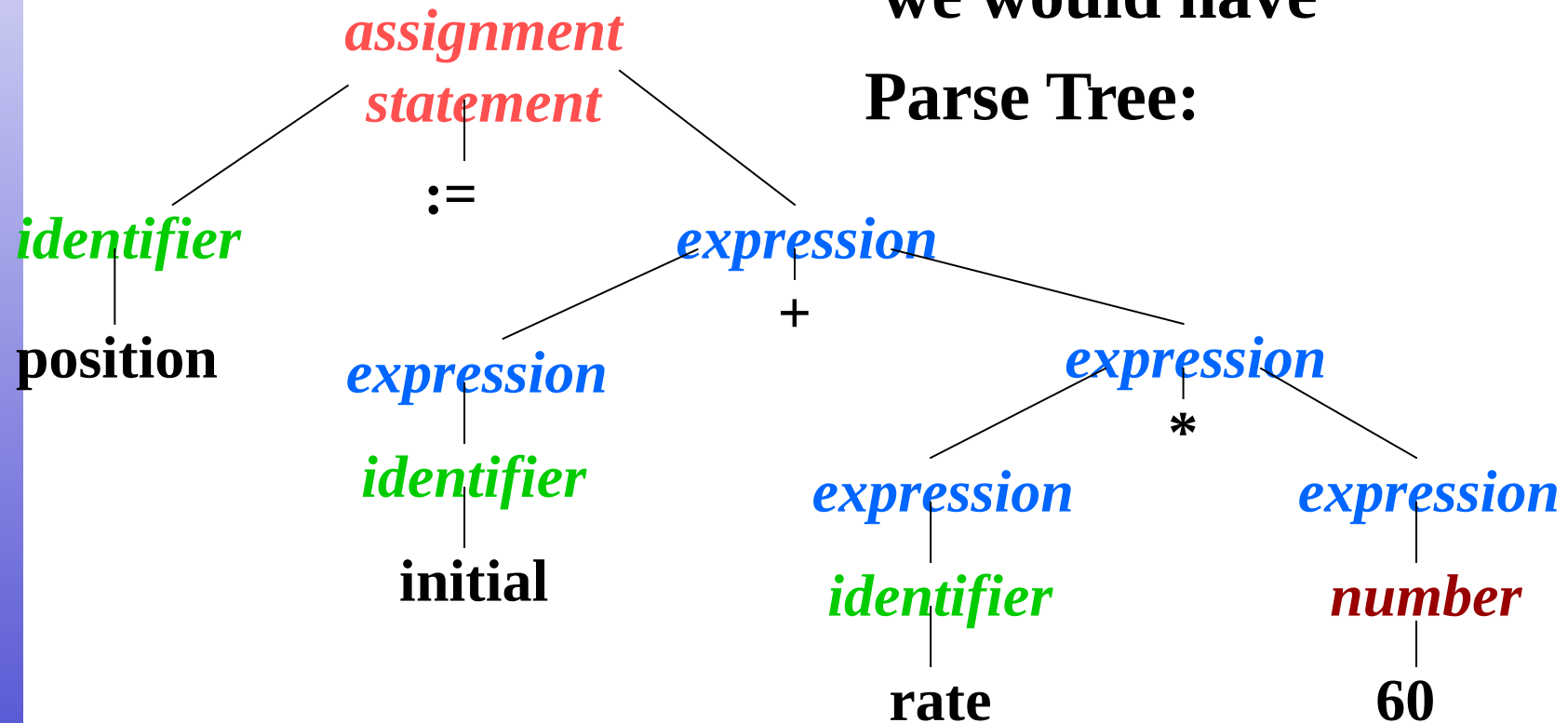
Example: **Position := initial + rate * 60 ;**



Blanks, Line breaks, etc. are scanned out

Phase 2. Syntax Analysis or Parsing

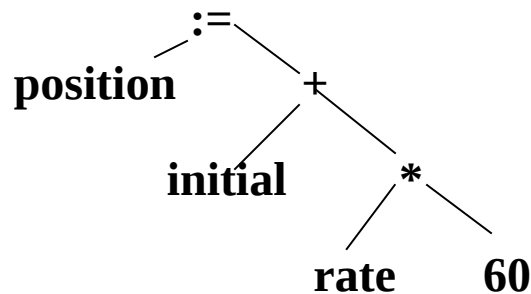
For previous example,
we would have
Parse Tree:



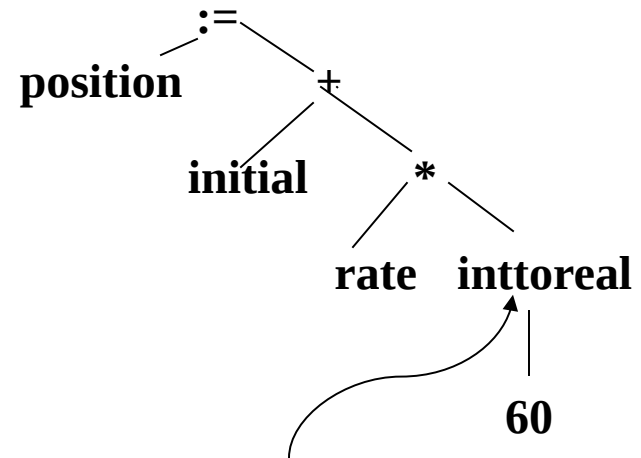
Nodes of tree are constructed using a grammar for the language

Phase 3. Semantic Analysis

- Finds Semantic Errors



Syntax Tree



Conversion Action

- One of the Most Important Activity in This Phase:
- Type Checking - Legality of Operands

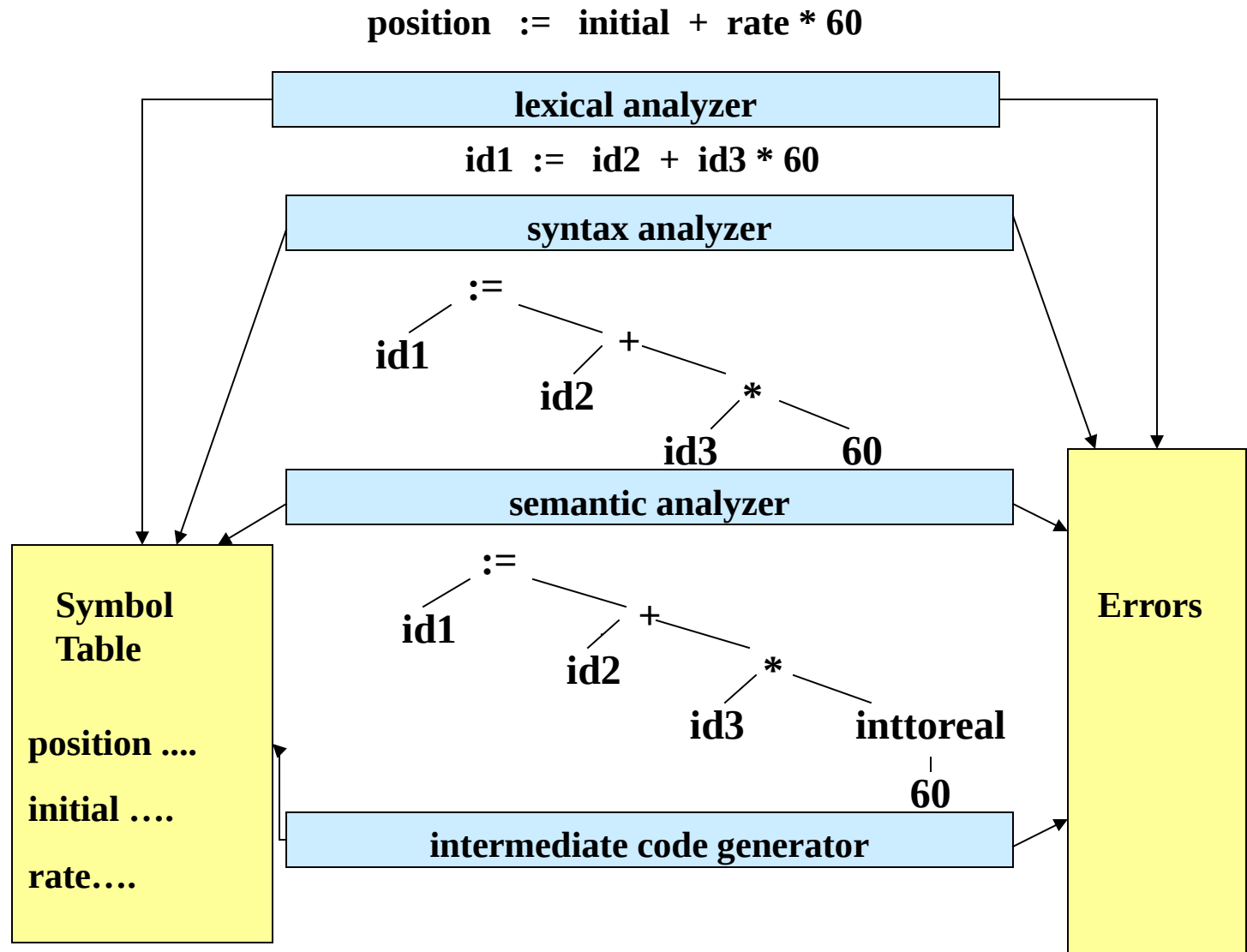
Supporting Phases/ Activities for Analysis

- Symbol Table Creation / Maintenance
 - Contains Info (storage, type, scope, args) on Each “Meaningful” Token, Typically Identifiers
 - Data Structure Created / Initialized During Lexical Analysis
 - Utilized / Updated During Later Analysis & Synthesis
- Error Handling
 - Detection of Different Errors Which Correspond to All Phases
 - What Happens When an Error Is Found?

The Synthesis Task For Compilation

- Intermediate Code Generation
 - Abstract Machine Version of Code - Independent of Architecture
 - Easy to Produce and Do Final, Machine Dependent Code Generation
- Code Optimization
 - Find More Efficient Ways to Execute Code
 - Replace Code With More Optimal Statements
- Final Code Generation
 - Generate Relocatable Machine Dependent Code
- Peephole Optimization
 - With a Very Limited View Improves Produced Final Code

Reviewing the Entire Process



Reviewing the Entire Process

